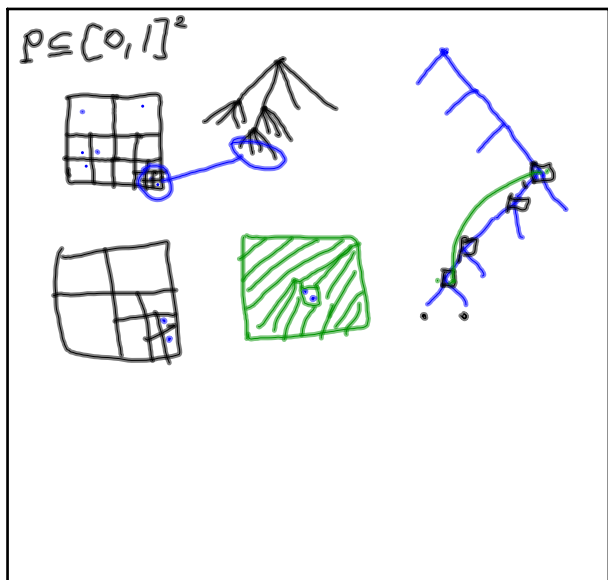
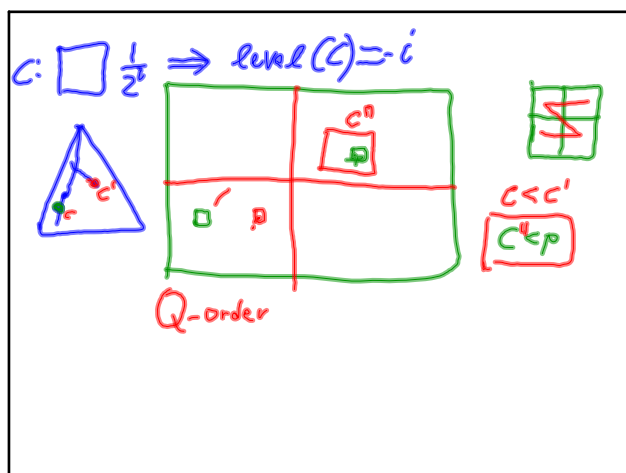




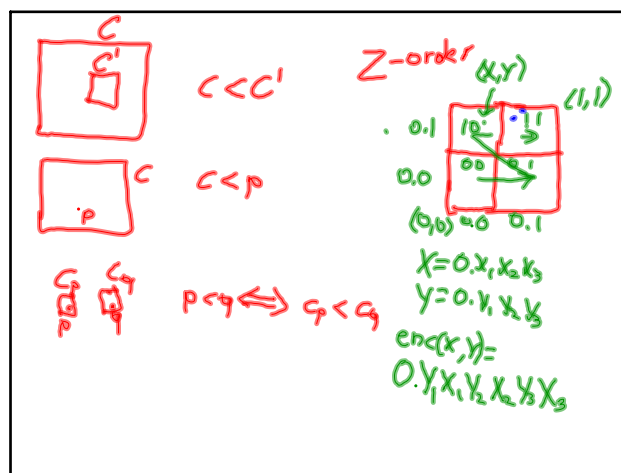
aug 18-08:59

$$\begin{aligned} \text{spread}(P) &= \frac{\dim(P)}{c_P(P)} \\ \text{QT for } P: \dim(P) &= \frac{1}{2} \\ \text{height}(\text{QT}(P)) &= \log \underbrace{\text{spread}(P)}_{\Phi} \\ \Rightarrow |\text{QT}(P)| &= O(n \log \underbrace{\Phi}_{\frac{1}{2}}) \\ \text{CQT: } |\text{CQT}(P)| &= O(n). \end{aligned}$$

aug 18-09:08



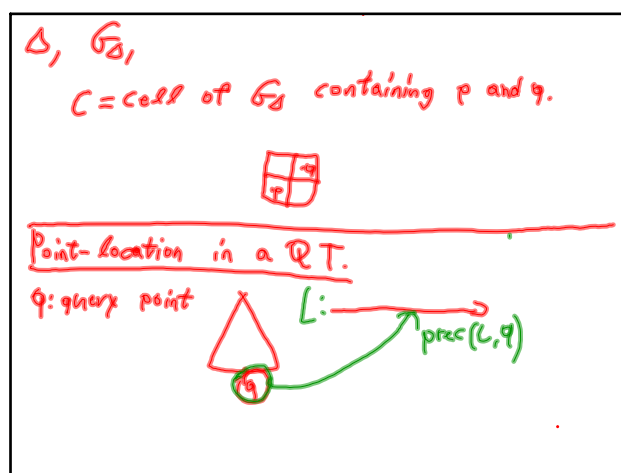
aug 18-09:13



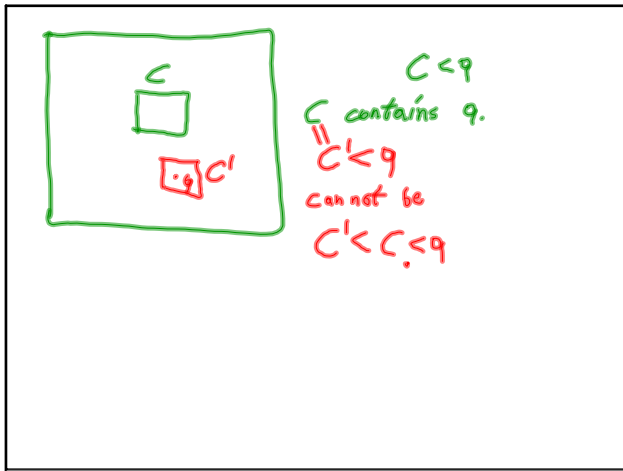
aug 18-09:20

$\text{bit}(x, y) = \text{index of first bit where numbers differ in base 2.}$
 $\text{bit}(\frac{3}{4}, \frac{7}{8}) = \text{bit}(0.11, 0.111) = 3$
 $p, q \quad \text{lca}(p, q) - \text{smallest common sq containing } p \text{ and } q$
 $\ell(\text{lca}(p, q)) = 1 - \min(\text{bit}(x_p, x_q), \text{bit}(y_p, y_q))$
 $\text{sidelen}(\text{lca}) = 2^{\text{level}(\text{lca})} = \Delta$

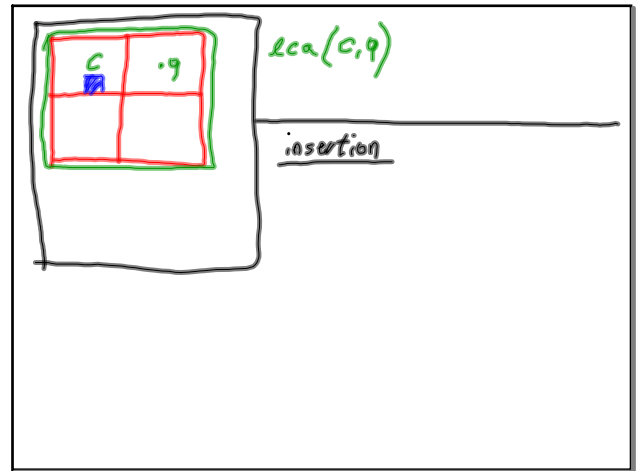
aug 18-09:31



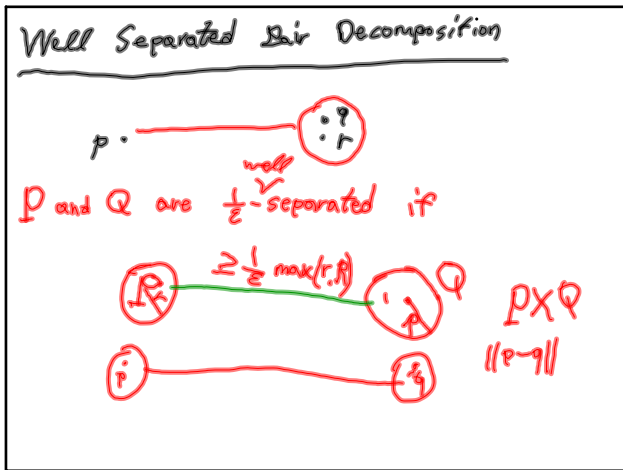
aug 18-09:40



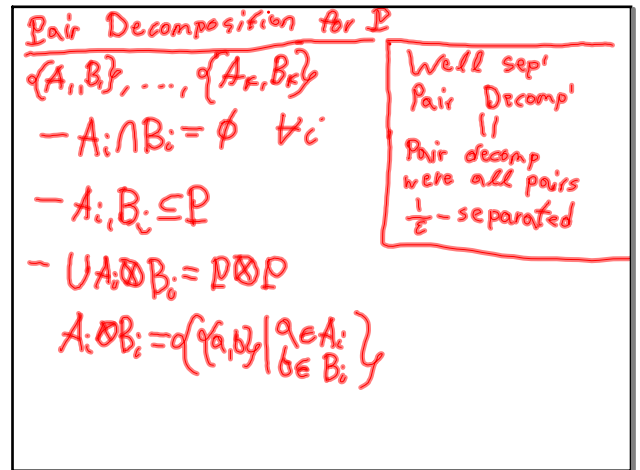
aug 18-09:45



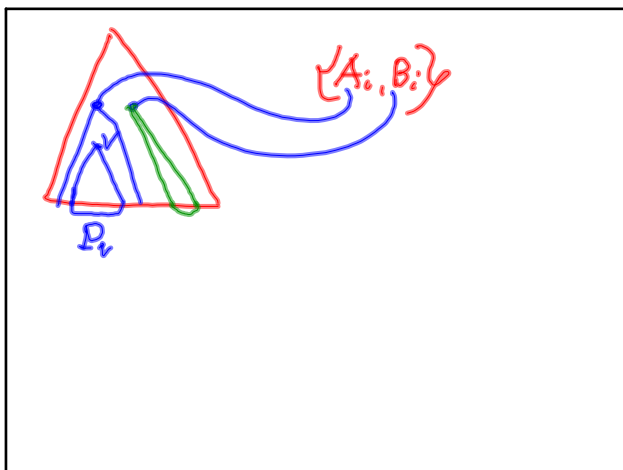
aug 18-09:48



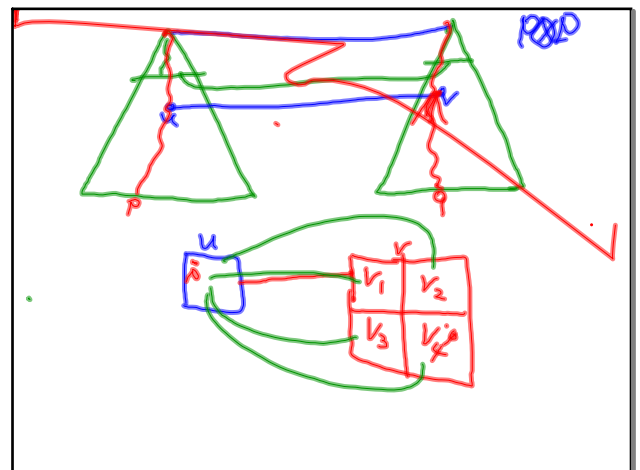
aug 18-09:56



aug 18-10:04



aug 18-10:11



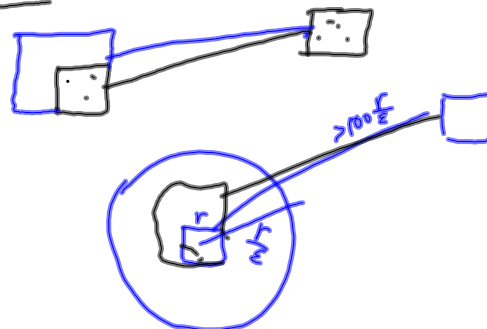
aug 18-10:13

Thm

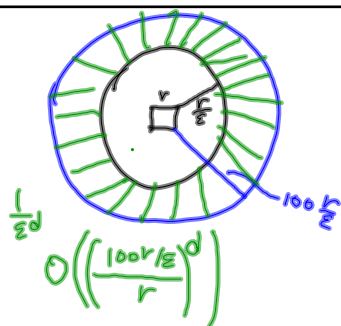
P : Set of n points in $\mathbb{R}^d$
 One can compute $\frac{1}{\epsilon}$ -WSPD of P
 of size $O(\frac{n}{\epsilon^d})$. RT is
 $O(n \log n + \frac{n}{\epsilon^d})$

aug 18-10:21

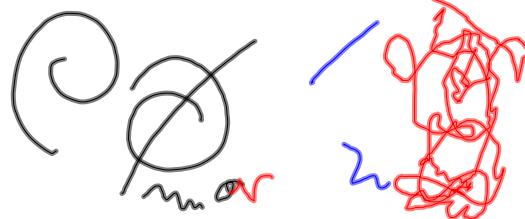
"Proof"



aug 18-10:24



aug 18-10:28



aug 18-10:34